Algorithms for Solving Higher Index DAEs

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Agenda

Part I
DAE Basics

Part II
Matching

Part III
BLT Sorting

Part IV
Pantelides

Part V
Dummy Derivatives

Part I

DAE Basics

DAE

System of Differential algebraic equation (DAE) in general form:

$$F(x, \dot{x}, y, t) = 0$$

where $x, \in \mathbb{R}^n, \dot{x} \in \mathbb{R}^n, y \in \mathbb{R}^m, F : G \subseteq \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R} \to \mathbb{R}^{n+m}$.

We have *n* number of variables that appear differentiated

We have *m* number of variables that *do not* appear differentiated

We have *n+m* number of equations

$$\dot{x} = -x$$
$$x(0) = 1$$

ODE, initial value problem (IVP)

$$\dot{x} = -x + y$$
$$x^2 + y^2 = 10$$

1 variable *x* is differentiated
1 variable y is not differentiated
2 equations

Is this an ODE or an DAE?

DAE, Example 1

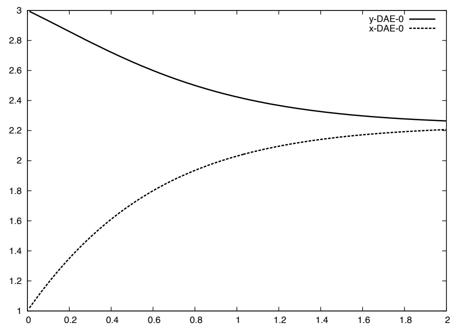
$$\dot{x} = -x + y$$
$$x^2 + y^2 = 10$$

Is it an initial value problem (IVP)?

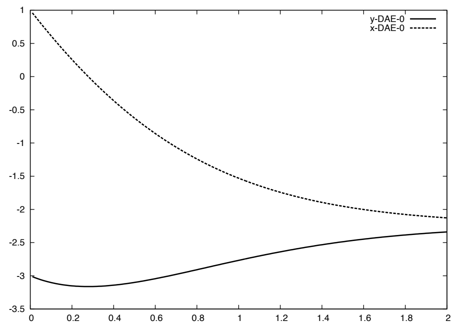
What should the initial value for y be?

$$x(0) = 1$$
$$y(0) = 3$$

We need to find consistent initial values.
Note that y(0) = -3 is also a consistent initial value.



$x(0) = 1 \quad y(0) = 3$



$x(0) = 1 \quad y(0) = -3$

DAE, Example 1

$$\dot{x} = -x + y$$
$$x^2 + y^2 = 10$$

Can we find an order that can solve these equations?

Yes, in each step:

1. Solve for *y* in equation (2). x is known.
2. Solve for x' in equation (1). Now both x and y are known.

In this case, we can actually symbolically transform this into an ODE directly.

$$\dot{x} = -x + \sqrt{10 - x^2}$$

(note that the DAE is nonlinear, and we here just selected one solution)

DAE, Example 2

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$$\dot{x} = -x + y - z$$
$$z = x^2 + y^2$$
$$z = x + x * y$$

Is this an DAE?

Yes, one differentiated variable (x) and two algebraic variables (y and z)

Can we find an order that can solve these equations?

No, equations 2 and 3 are algebraically dependent on each other.

$$\dot{x} = f(x, y, t)$$
$$0 = g(x, y, t)$$

This is called the *semi-explicit form* of an DAE

Solution (in each time step)

1. Solve (nonlinear) algebraic equations
2. Solve differentiated variables

DAE Index

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<u>Definition</u>: The index of an DAE is the minimum number of times that all or part of the DAE must be differentiated with respect to *t* in order to determine *x'* as a continuous function of *x* and *t*.

(Brenan, Campbell, Petzold, 1989)

$$\dot{x} = -x + y$$
$$x^2 + y^2 = 10$$

Our first example was an index 0 DAE. No differentiation is need to obtain an ODE. An ODE has also index 0.

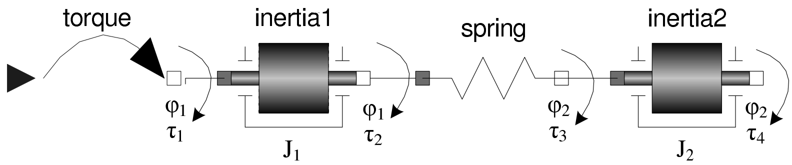
$$\dot{x} = -x + y - z$$
$$z = x^2 + y^2$$
$$z = x + x * y$$

Example two has an algebraic loop, and the two algebraic equations are non-singular. Example of an index 1 DAE.

Note that you can differentiate parts of the equation system once (equations (2) and (3)) to obtain an ODE. (Not recommended for numerical stability)

We will soon see examples where a system of equation is singular. These may be *higher-index DAEs (index > 1).*

Drive Shaft Example



Is this an DAE?

$$
\begin{aligned}
\dot{\varphi}_1 &= \omega_1 \\
\dot{\varphi}_2 &= \omega_2 \\
\dot{\omega}_1 &= \frac{\tau_1 + \tau_2}{J_1} \\
\dot{\omega}_2 &= \frac{\tau_3 + \tau_4}{J_2} \\
\tau_1 &= u \\
\tau_2 &= c \cdot (\varphi_2 - \varphi_1) \\
\tau_3 &= -c \cdot (\varphi_2 - \varphi_1) \\
\tau_4 &= 0
\end{aligned}
$$

Variables: $(\varphi_1, \varphi_2, \omega_1, \omega_2, \tau_1, \tau_2, \tau_3, \tau_4)$

Appearing differentiated: $(\dot{\varphi}_1, \dot{\varphi}_2, \dot{\omega}_1, \dot{\omega}_2)$

Incidence matrix. Differentiated variables and the algebraic variables are unknown.

	$\dot{\omega}_1$	$\dot{\omega}_2$	$\dot{\phi}_1$	$\dot{\phi}_2$	τ_1	τ_2	τ_3	τ_4
f_1	0	0	1	0	0	0	0	0
f_2	0	0	0	1	0	0	0	0
f_3	1	0	0	0	1	1	0	0
f_4	0	1	0	0	0	0	1	1
f_5	0	0	0	0	1	0	0	0
f_6	0	0	0	0	0	1	0	0
f_7	0	0	0	0	0	0	1	0
f_8	0	0	0	0	0	0	0	1

	τ_4	τ_3	τ_2	τ_1	$\dot{\omega}_2$	$\dot{\omega}_1$	$\dot{\phi}_2$	$\dot{\phi}_1$
f_8	1	0	0	0	0	0	0	0
f_7	0	1	0	0	0	0	0	0
f_6	0	0	1	0	0	0	0	0
f_5	0	0	0	1	0	0	0	0
f_4	1	1	0	0	1	0	0	0
f_3	0	0	1	1	0	1	0	0
f_2	0	0	0	0	0	0	1	0
f_1	0	0	0	0	0	0	0	1

Index-0 DAE
(by substitution we get directly an ODE)

Matching: Find a unique mapping between variables and equations.

Sorting: Sort equations (permute matrix)

Part II

Matching

Example: Matching

System of equations

$$f_1(y) = 0$$
$$f_2(\dot{x}_1, \dot{x}_2, y) = 0$$
$$f_3(\dot{x}_2) = 0$$

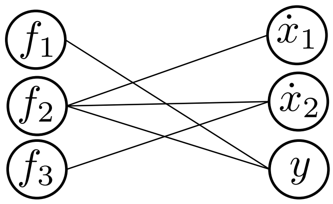
Construct a bipartite graph

$$G = (F, V, E)$$

$$F = \{f_1, f_2, f_3\} \quad E = \{(f_1, y), (f_2, \dot{x}_1), (f_2, \dot{x}_2), (f_2, y), (f_3, \dot{x}_2)\}$$
$$V = \{\dot{x}_1, \dot{x}_2, y\}$$

Incidence Matrix

	$\dot{x}_1$	$\dot{x}_2$	y
f_1	0	0	1
f_2	1	1	1
f_3	0	1	0



Algorithm: Matching


```
MATCH(G)
  assign ← ∅
  for each f ∈ G.F
      do C ← ∅
         if not MATCH-EQUATION(G, f, C, assign, ∅)
            then return (FALSE, assign)
  return (TRUE, assign)
```

Color visited vertices

$$C \subseteq G.F \cup G.V$$

Assigns variables to equations

$$assign[v] = \begin{cases} f & \text{if } f \text{ matches } v \\ \text{NIL} & \text{otherwise} \end{cases}$$

Underline means call by reference.

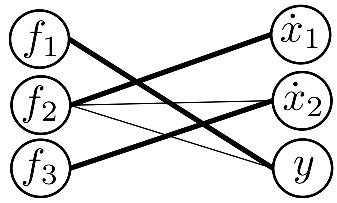
vmap and *equation coloring* is not used until in Part IV.

```
MATCH-EQUATION(G, f, C, assign, vmap)
  C ← C ∪ {f}
  if there exits a v ∈ G.V such that (f, v) ∈ G.E
     and assign[v] = NIL and vmap[v] = NIL
    then assign[v] ← f
         return TRUE
    else for each v where (f, v) ∈ G.E and v ∉ C
           and vmap[v] = NIL
           do C ← C ∪ {v}
              if MATCH-EQUATION(G, assign[v], C, assign, vmap)
                then assign[v] ← f
                     return TRUE
  return FALSE
```

Example: Matching

```
MATCH(G)
  assign ← ∅
  for each f ∈ G.F
      do C ← ∅
          if not MATCH-EQUATION(G, f, C, assign, ∅)
             then return (FALSE, assign)
  return (TRUE, assign)
```

```
MATCH-EQUATION(G, f, C, assign, vmap)
  C ← C ∪ {f}
  if there exits a v ∈ G.V such that (f, v) ∈ G.E
     and assign[v] = NIL and vmap[v] = NIL
    then assign[v] ← f
         return TRUE
    else for each v where (f, v) ∈ G.E and v ∉ C
           and vmap[v] = NIL
             do C ← C ∪ {v}
                if MATCH-EQUATION(G, assign[v], C, assign, vmap)
                  then assign[v] ← f
                       return TRUE
  return FALSE
```



Exercise

Do each step of the algorithms and keep track of *C* and *assign*.

Case A: For f1, use x1.

$assign = \{y \mapsto f_1, \dot{x}_1 \mapsto f_2, \dot{x}_2 \mapsto f_3\}$

$C = \{f_1, f_2, f_3\}$

Case B: For f1, first use x2 (Reassignment of x2)

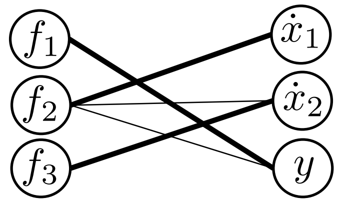
$assign = \{y \mapsto f_1, \dot{x}_1 \mapsto f_2, \dot{x}_1 \mapsto f_3\}$

$C = \{f_1, f_2, f_3, \dot{x}_2, \}$

Example: Matching

System of equations

$$f_1(y) = 0$$
$$f_2(\dot{x}_1, \dot{x}_2, y) = 0$$
$$f_3(\dot{x}_2) = 0$$



Incidence Matrix

$$\begin{array}{c} \\ f_1 \\ f_2 \\ f_3 \end{array} \begin{array}{c} \begin{array}{ccc} \dot{x}_1 & \dot{x}_2 & y \end{array} \\ \begin{pmatrix} 0 & 0 & \boxed{1} \\ \boxed{1} & 1 & 1 \\ 0 & \boxed{1} & 0 \end{pmatrix} \end{array}$$

We may now permute the matrix

$$\begin{array}{c} \\ f_3 \\ f_1 \\ f_2 \end{array} \begin{array}{c} \begin{array}{ccc} \dot{x}_2 & y & \dot{x}_1 \end{array} \\ \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 1 & 1 \end{pmatrix} \end{array}$$

The matching problem solves the problem of finding a permutation such that the matrix has a nonzero diagonal. Also called *maximum traversal*.

Sorting into Lower Triangular Form

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	x_1	x_2	x_3	x_4	x_5	x_6	x_7
f_1	0	0	1	0	0	1	0
f_2	1	1	0	0	1	1	0
f_3	0	0	0	1	1	1	0
f_4	0	1	0	0	0	0	0
f_5	0	0	1	0	1	0	0
f_6	0	1	0	0	0	1	0
f_7	0	1	1	0	0	0	1

Unsorted Matrix

Find a matching

	x_2	x_6	x_3	x_7	x_5	x_4	x_1
f_4	1	0	0	0	0	0	0
f_6	1	1	0	0	0	0	0
f_1	0	1	1	0	0	0	0
f_7	1	0	1	1	0	0	0
f_5	0	0	1	0	1	0	0
f_3	0	1	0	0	1	1	0
f_2	1	1	0	0	1	0	1

Sorting (permutation of Matrix) into <u>Lower Triangular Matrix Form</u>

We now have a causal form; solving the equation system is straight forward.

But, we cannot always permute the matrix into lower triangular form…

An DAE (or ODE) in Lower triangular matrix form is index 1.

Sorting into Block Lower Triangular (BLT) Form

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	x_1	x_2	x_3	x_4	x_5	x_6	x_7
f_1	0	0	1	1	0	1	0
f_2	1	1	0	0	1	1	0
f_3	0	0	0	1	1	0	0
f_4	0	1	0	0	0	0	0
f_5	1	0	1	0	1	0	0
f_6	0	1	0	0	0	1	0
f_7	0	0	1	0	0	0	1

Another unsorted Matrix

	x_2	x_6	x_3	x_5	x_1	x_4	x_7
f_4	1	0	0	0	0	0	0
f_6	1	1	0	0	0	0	0
f_1	0	1	1	0	0	1	0
f_5	0	0	1	1	1	0	0
f_2	1	1	0	1	1	0	0
f_3	0	0	0	1	0	1	0
f_7	0	0	1	0	0	0	1

Sorting (permutation of Matrix) into Block <u>Lower Triangular (BLT) Form</u>

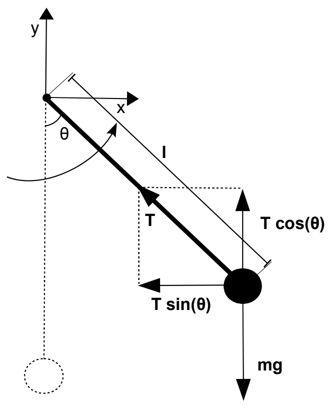
We have identified an <u>**algebraic loop**</u>.

At each time step, the algebraic loops may be solved using Guassian elimination (if linear) or a Newton's method (if nonlinear).

In part III we discuss a BLT sorting algorithm

An DAE in BLT form with algebraic loops (structurally non-singular) is Index 1

Example: Pendulum



Pendulum in Cartesian coordinate system

$$-T \cdot \frac{x}{l} = m\ddot{x}$$
$$-T \cdot \frac{y}{l} - mg = m\ddot{y}$$
$$x^2 + y^2 = l^2$$

Simplified using

$$-T/l = \lambda$$
$$m = 1$$
$$l^2 = L$$

Simplified

$$\ddot{x} = \lambda \cdot x$$
$$\ddot{y} = \lambda \cdot y - g$$
$$x^2 + y^2 = L$$

Rewritten in first order

$$\dot{x} = u$$
$$\dot{y} = v$$
$$\dot{u} = \lambda \cdot x$$
$$\dot{v} = \lambda \cdot y - g$$
$$x^2 + y^2 = L$$

Incidence Matrix

$$\begin{array}{c} \\ f_1 \\ f_2 \\ f_3 \\ f_4 \\ f_5 \end{array} \begin{array}{c} \begin{array}{ccccc} \dot{x} & \dot{y} & \dot{u} & \dot{v} & \lambda \end{array} \\ \begin{pmatrix} \mathbf{1} & 0 & 0 & 0 & 0 \\ 0 & \mathbf{1} & 0 & 0 & 0 \\ 0 & 0 & \mathbf{1} & 0 & \mathbf{1} \\ 0 & 0 & 0 & \mathbf{1} & \mathbf{1} \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix} \end{array}$$

Is this an DAE?

Can we solve it?

Can we create BLT?

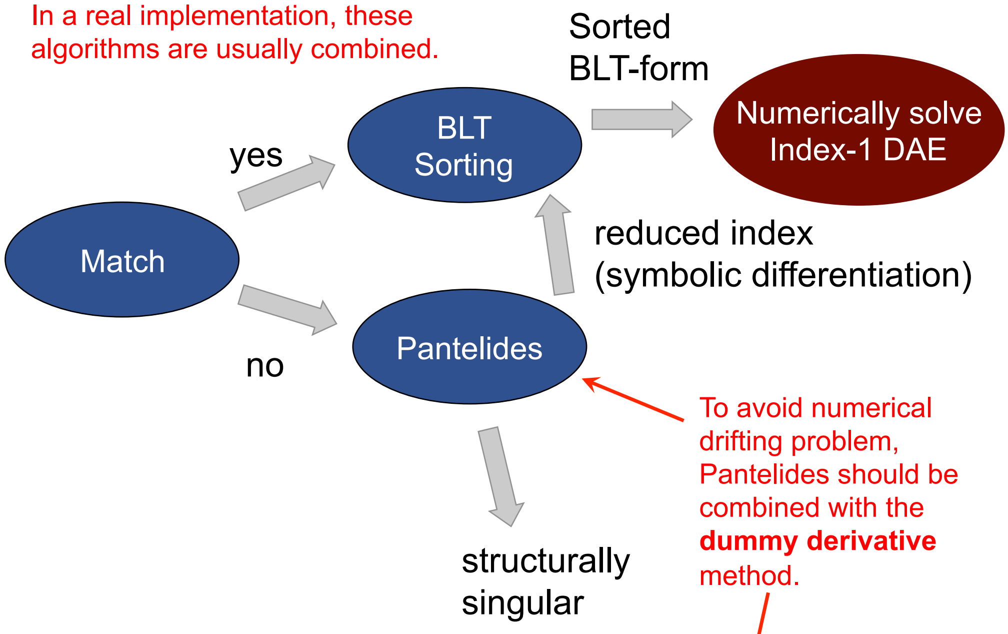
IDEA: Symbolically differentiate equations to get derivatives.

No, we cannot find a matching (see f_5). This is a higher-index problem (index > 1).

How should we determine which equations to differentiate. Solution: Pantelides Algorithm

How the algorithms fit together (simplified)

Note that this is a **simplified view**: In a real implementation, these algorithms are usually combined.



Part III

BLT Sorting

Algorithm: BLT Sort

Input: a bipartite graph G

```
BLT(G)
  (match, assign) ← MATCH(G)
  if not match
     then return error "Singular"

  D.V ← G.F
  D.E ← ∅
  for each (f, v) ∈ G.E where f ∈ G.F and assign[v] ≠ f
      do D.E ← D.E ∪ {(assign[v], f)}

  MAKEEMPTY(O)
  MAKEEMPTY(S)
  i ← 0
  lowlink ← ∅
  number ← ∅
  for each v ∈ D.V
      do if number[v] = NIL
            then STRONGCONNECT(v, D, S, i, lowlink, number, O)
  return O
```

Output: a stack of sets of equation vertices, where each set represents an equation block in the BLT matrix.

Algorithm: BLT Sort

Input: a bipartite graph G

BLT(G)

Part 1
Find matching

```
(match, assign) ← MATCH(G)
if not match
   then return error "Singular"
```

Part 2
Construct equation dependency graph

```
D.V ← G.F
D.E ← ∅
for each (f, v) ∈ G.E where f ∈ G.F and assign[v] ≠ f
    do D.E ← D.E ∪ {(assign[v], f)}
```

Part 3
Sort into blocks of equations using Tarjan's strongly connected component algorithm

```
MAKEEMPTY(O)
MAKEEMPTY(S)
i ← 0
lowlink ← ∅
number ← ∅
for each v ∈ D.V
    do if number[v] = NIL
          then STRONGCONNECT(v, D, S, i, lowlink, number, O)
return O
```

Output: a stack of sets of equation vertices, where each set represents an equation block in the BLT matrix.

Example: BLT Sort

$$
\begin{array}{c} \\ f_1 \\ f_2 \\ f_3 \\ f_4 \\ f_5 \\ f_6 \end{array}
\begin{array}{c}
\begin{array}{cccccc} \dot{x}_1 & \dot{x}_2 & \dot{x}_3 & \dot{x}_4 & y_1 & y_2 \end{array} \\
\begin{pmatrix}
0 & 0 & \mathbf{1} & \boxed{1} & 0 & \mathbf{1} \\
\boxed{1} & \mathbf{1} & 0 & 0 & \mathbf{1} & \mathbf{1} \\
0 & \boxed{1} & \mathbf{1} & 0 & 0 & 0 \\
\mathbf{1} & 0 & \mathbf{1} & 0 & \boxed{1} & 0 \\
0 & \mathbf{1} & 0 & 0 & 0 & \boxed{1} \\
0 & 0 & \boxed{1} & 0 & 0 & \mathbf{1}
\end{pmatrix}
\end{array}
$$

$$G = (F, V, E)$$

$$F = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$

$$V = \{\dot{x}_1, \dot{x}_2, \dot{x}_3, \dot{x}_4, y_1, y_2\}$$

In Part 1 of BLT - matching

```
1  (match, assign) ← MATCH(G)
2  if not match
3     then return error "Singular"
```

Returns TRUE (steps omitted) with assignment

$$assign = \{\dot{x}_1 \mapsto f_2, \dot{x}_2 \mapsto f_3, \dot{x}_3 \mapsto f_6, \dot{x}_4 \mapsto f_1, y_1 \mapsto f_4, y_2 \mapsto f_5\}$$

Algorithm: BLT Sort

Input: a bipartite graph G

```
BLT(G)
(match, assign) ← MATCH(G)
if not match
   then return error "Singular"

D.V ← G.F
D.E ← ∅
for each (f, v) ∈ G.E where f ∈ G.F and assign[v] ≠ f
    do D.E ← D.E ∪ {(assign[v], f)}

MAKEEMPTY(O)
MAKEEMPTY(S)
i ← 0
lowlink ← ∅
number ← ∅
for each v ∈ D.V
    do if number[v] = NIL
          then STRONGCONNECT(v, D, S, i, lowlink, number, O)
return O
```

Part 2
Construct equation dependency graph

Output: a stack of sets of equation vertices, where each set represents an equation block in the BLT matrix.

Example: BLT Sort

$$\begin{array}{c} \\ f_1 \\ f_2 \\ f_3 \\ f_4 \\ f_5 \\ f_6 \end{array} \begin{array}{c} \begin{array}{cccccc} \dot{x}_1 & \dot{x}_2 & \dot{x}_3 & \dot{x}_4 & y_1 & y_2 \end{array} \\ \begin{pmatrix} 0 & 0 & 1 & \boxed{1} & 0 & 1 \\ \boxed{1} & 1 & 0 & 0 & 1 & 1 \\ 0 & \boxed{1} & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & \boxed{1} & 0 \\ 0 & 1 & 0 & 0 & 0 & \boxed{1} \\ 0 & 0 & \boxed{1} & 0 & 0 & 1 \end{pmatrix} \end{array}$$

When solving x'$_4$, variables x'$_3$ and y_2 need to be known. Hence, f_1 is dependent on f_5 and f_6.

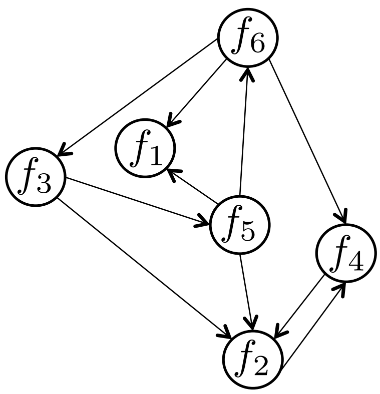
$$G = (F, V, E) \qquad F = \{f_1, f_2, f_3, f_4, f_5, f_6\} \quad V = \{\dot{x}_1, \dot{x}_2, \dot{x}_3, \dot{x}_4, y_1, y_2\}$$

Matching tells us that x'$_3$ is supposed to be solved using f_6

$$assign = \{\dot{x}_1 \mapsto f_2, \dot{x}_2 \mapsto f_3, \dot{x}_3 \mapsto f_6, \dot{x}_4 \mapsto f_1, y_1 \mapsto f_4, y_2 \mapsto f_5\}$$

In Part 2 of BLT – construct equation dependency graph (digraph)

```
D.V ← G.F
D.E ← ∅
for each (f, v) ∈ G.E where f ∈ G.F and assign[v] ≠ f
    do D.E ← D.E ∪ {(assign[v], f)}
```



$$D = (V, E)$$

$$V = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$

$$E = \{f_2 \mapsto f_4, f_3 \mapsto f_2, f_3 \mapsto f_5, f_4 \mapsto f_2, f_5 \mapsto f_1, f_5 \mapsto f_2, f_5 \mapsto f_6, f_6 \mapsto f_1, f_6 \mapsto f_3, f_6, \mapsto f_4\}$$

Exercise
Create D graphically

Algorithm: BLT Sort

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Input: a bipartite graph G

```
BLT(G)
(match, assign) ← MATCH(G)
if not match
   then return error "Singular"

D.V ← G.F
D.E ← ∅
for each (f, v) ∈ G.E where f ∈ G.F and assign[v] ≠ f
    do D.E ← D.E ∪ {(assign[v], f)}

MAKEEMPTY(O)
MAKEEMPTY(S)
i ← 0
lowlink ← ∅
number ← ∅
for each v ∈ D.V
    do if number[v] = NIL
          then STRONGCONNECT(v, D, S, i, lowlink, number, O)
return O
```

Part 3
Sort into blocks of equations using Tarjan's strongly connected component algorithm

Output: a stack of sets of equation vertices, where each set represents an equation block in the BLT matrix.

Algorithm: StrongConnect (Tarjan)

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```
STRONGCONNECT(v, D, S, i, lowlink, number, O)
i ← i + 1
lowlink[v] ← i
number[v] ← i
PUSH(S, v)
for each w ∈ D.V where (v, w) ∈ D.E
    do if number[w] = NIL
          then STRONGCONNECT(w, D, S, i, lowlink, number, O)
               lowlink[v] ← MIN(lowlink[v], lowlink[w])
          else if w ∈ S and number[w] < number[v]
                  then lowlink[v] ← MIN(lowlink[v], number[w])
if lowlink[v] = number[v]
  then eqset ← ∅
       while not ISEMPTY(S) and number[TOP(S)] ≥ number[v]
          do eqset ← eqset ∪ {POP(S)}
       PUSH(O, eqset)
return
```

Algorithm: StrongConnect (Tarjan)

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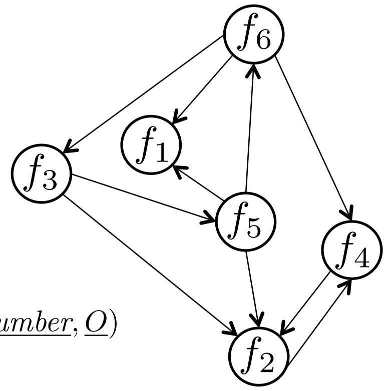


```
STRONGCONNECT(v, D, S, i, lowlink, number, O)
i ← i + 1
lowlink[v] ← i
number[v] ← i
PUSH(S, v)
for each w ∈ D.V where (v, w) ∈ D.E
    do if number[w] = NIL
         then STRONGCONNECT(w, D, S, i, lowlink, number, O)
              lowlink[v] ← MIN(lowlink[v], lowlink[w])
         else if w ∈ S and number[w] < number[v]
                then lowlink[v] ← MIN(lowlink[v], number[w])
if lowlink[v] = number[v]
   then eqset ← ∅
        while not ISEMPTY(S) and number[TOP(S)] ≥ number[v]
           do eqset ← eqset ∪ {POP(S)}
        PUSH(O, eqset)
return
```

Exercise
Construct stack O

Top of the stack is to the left

$O = [\{f_3, f_5, f_6\}, \{f_2, f_4\}, \{f_1\}]$

Example: BLT Sort

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	$\dot{x}_1$	$\dot{x}_2$	$\dot{x}_3$	$\dot{x}_4$	y_1	y_2
f_1	0	0	1	[1]	0	1
f_2	[1]	1	0	0	1	1
f_3	0	[1]	1	0	0	0
f_4	1	0	1	0	[1]	0
f_5	0	1	0	0	0	[1]
f_6	0	0	[1]	0	0	1

$$G = (F, V, E)$$

$$F = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$

$$V = \{\dot{x}_1, \dot{x}_2, \dot{x}_3, \dot{x}_4, y_1, y_2\}$$

$$assign = \{\dot{x}_1 \mapsto f_2, \dot{x}_2 \mapsto f_3, \dot{x}_3 \mapsto f_6, \dot{x}_4 \mapsto f_1, y_1 \mapsto f_4, y_2 \mapsto f_5\}$$

$$O = [\{f_3, f_5, f_6\}, \{f_2, f_4\}, \{f_1\}]$$

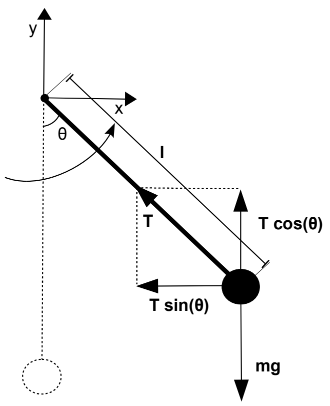
We can now create the sorted BLT matrix

	$\dot{x}_2$	y_2	$\dot{x}_3$	$\dot{x}_1$	y_1	$\dot{x}_4$
f_3	1	0	1	0	0	0
f_5	1	1	0	0	0	0
f_6	0	1	1	0	0	0
f_2	1	1	0	1	1	0
f_4	0	0	1	1	1	0
f_1	0	1	1	0	0	1

Part IV

Pantelides

Example: Pendulum



Pendulum in Cartesian coordinate system

$$-T \cdot \frac{x}{l} = m\ddot{x}$$
$$-T \cdot \frac{y}{l} - mg = m\ddot{y}$$
$$x^2 + y^2 = l^2$$

Simplified using

$$-T/l = \lambda$$
$$m = 1$$
$$l^2 = L$$

Simplified

$$\ddot{x} = \lambda \cdot x$$
$$\ddot{y} = \lambda \cdot y - g$$
$$x^2 + y^2 = L$$

Rewritten in first order

$$\dot{x} = u$$
$$\dot{y} = v$$
$$\dot{u} = \lambda \cdot x$$
$$\dot{v} = \lambda \cdot y - g$$
$$x^2 + y^2 = L$$

Incidence Matrix

$$\begin{array}{c} \\ f_1 \\ f_2 \\ f_3 \\ f_4 \\ f_5 \end{array} \begin{array}{c} \begin{array}{ccccc} \dot{x} & \dot{y} & \dot{u} & \dot{v} & \lambda \end{array} \\ \begin{pmatrix} \mathbf{1} & 0 & 0 & 0 & 0 \\ 0 & \mathbf{1} & 0 & 0 & 0 \\ 0 & 0 & \mathbf{1} & 0 & \mathbf{1} \\ 0 & 0 & 0 & \mathbf{1} & \mathbf{1} \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix} \end{array}$$

Pendulum: Graph Construction

System of equations

$$\dot{x} = u$$
$$\dot{y} = v$$
$$\dot{u} = \lambda \cdot x$$
$$\dot{v} = \lambda \cdot y - g$$
$$x^2 + y^2 = L$$

$$f_1(\dot{x}, u) = 0$$
$$f_2(\dot{y}, v) = 0$$
$$f_3(\dot{u}, \lambda, x) = 0$$
$$f_4(\dot{v}, \lambda, y) = 0$$
$$f_5(x, y) = 0$$

Note that we include both differentiated and not differentiated variables.

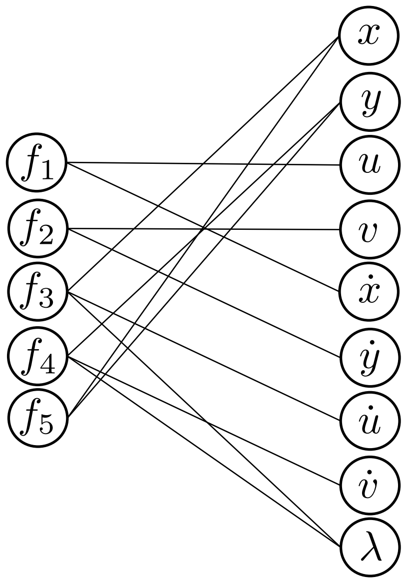
Construct a bipartite graph

$G = (F, V, E)$

$$F = \{f_1, f_2, f_3, f_4, f_5\}$$
$$V = \{x, y, u, v, \dot{x}, \dot{y}, \dot{u}, \dot{v}, \lambda\}$$
$$E = \{(f_1, \dot{x}), (f_1, u), (f_2, \dot{y}), (f_2, v), (f_3, \dot{u}), (f_3, \lambda), (f_3, x), (f_4, \dot{v}), (f_4, \lambda), (f_4, y), (f_5, x), (f_5, y)\}$$

Pendulum

$$f_1(\dot{x}, u) = 0$$
$$f_2(\dot{y}, v) = 0$$
$$f_3(\dot{u}, \lambda, x) = 0$$
$$f_4(\dot{v}, \lambda, y) = 0$$
$$f_5(x, y) = 0$$



Algorithm: Pantelides

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```
PANTELIDES(G, vmap, eqmap)
1   assign ← ∅
2   for each e ∈ G.F
3       do f ← e
4           repeat
5               C ← ∅
6               match ← MATCH-EQUATION(G, f, C, assign, vmap)
7               if not match
8                 then for each v ∈ C where v ∈ G.V
9                          do let v′ be a vertex, such that v′ ∉ G.V
10                            vmap[v] ← v′
11                            G.V ← G.V ∪ {v′}
12                      for each f ∈ C where f ∈ G.F
13                          do let f′ be a vertex, such that f′ ∉ G.F
14                             eqmap[f] ← f′
15                             G.F ← G.F ∪ {f′}
16                             for each v ∈ G.V where (f, v) ∈ G.E
17                                 do G.E ← G.E ∪ {(f′, v), (f′, vmap[v])}
18                      for each v ∈ C where v ∈ G.V
19                          do assign[vmap[v]] ← eqmap[assign[v]]
20                      f ← eqmap[f]
21          until match
22  return assign
```

Mapping variables to differentiated variables

$$vmap[v] = \begin{cases} v' & \text{if } \frac{dv}{dt} = v' \\ \text{NIL} & \text{otherwise} \end{cases}$$

Mapping equations to their differentiated version

$$eqmap[f] = \begin{cases} f' & \text{if } \frac{df}{dt} = f' \\ \text{NIL} & \text{otherwise} \end{cases}$$

Assigns variables to equations

$$assign[v] = \begin{cases} f & \text{if } f \text{ matches } v \\ \text{NIL} & \text{otherwise} \end{cases}$$

Pendulum

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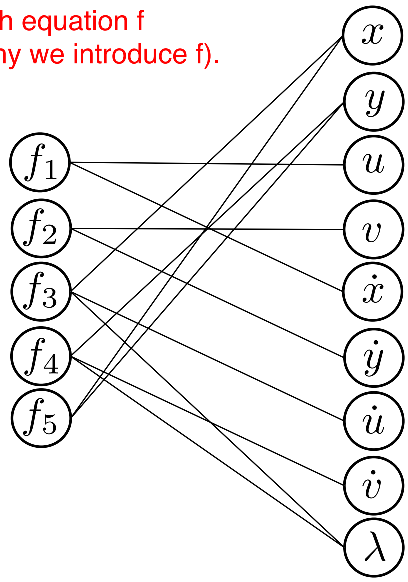
We start with no variable to equation assignments.

Iterate over each equation f (we see later why we introduce f).

```
PANTELIDES(G, vmap, eqmap)
1   assign ← ∅
2   for each e ∈ G.F
3       do f ← e
4           repeat
5               C ← ∅
```

Preparation for matching algorithm. Set all vertices to be uncolored.



Initial state after step 5.

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$$
$$eqmap = \{\}$$
$$assign = \{\}$$
$$C = \{\}$$

Algorithm: Pantelides

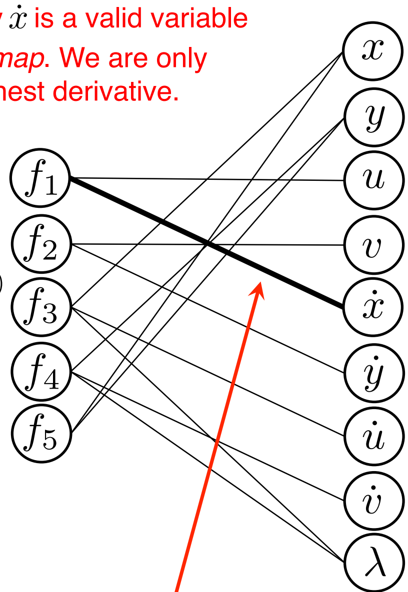
```
PANTELIDES(G, vmap, eqmap)
assign ← ∅
for each e ∈ G.F
    do f ← e
       repeat
              C ← ∅
              match ← MATCH-EQUATION(G, f, C, assign, vmap)
              if not match
                then for each v ∈ C where v ∈ G.V
                          do let v′ be a vertex, such that v′ ∉ G.V
                             vmap[v] ← v′
                             G.V ← G.V ∪ {v′}
                     for each f ∈ C where f ∈ G.F
                          do let f′ be a vertex, such that f′ ∉ G.F
                             eqmap[f] ← f′
                             G.F ← G.F ∪ {f′}
                             for each v ∈ G.V where (f, v) ∈ G.E
                                 do G.E ← G.E ∪ {(f′, v), (f′, vmap[v])}
                     for each v ∈ C where v ∈ G.V
                          do assign[vmap[v]] ← eqmap[assign[v]]
                     f ← eqmap[f]
         until match
return assign
```

Try to find a match for equation ***f***.

Pendulum

```
MATCH-EQUATION(G, f, C, assign, vmap)
C ← C ∪ {f}
if there exits a v ∈ G.V such that (f, v) ∈ G.E
   and assign[v] = NIL and vmap[v] = NIL
  then assign[v] ← f
       return TRUE
  else for each v where (f, v) ∈ G.E and v ∉ C
          and vmap[v] = NIL
          do C ← C ∪ {v}
             if MATCH-EQUATION(G, assign[v], C, assign, vmap)
               then assign[v] ← f
                    return TRUE
return FALSE
```

Note that only $\dot{x}$ is a valid variable because of *vmap*. We are only match for highest derivative.



Matched variable to equation

State when returning from Match-Equation.

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$$
$$eqmap = \{\}$$
$$assign = \{\dot{x} \mapsto f_1\}$$
$$C = \{f_1\}$$

Matched variable to equation

Colored one equation

Algorithm: Pantelides

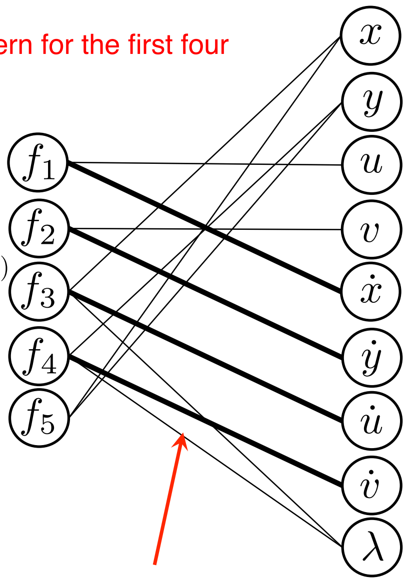
```
PANTELIDES(G, vmap, eqmap)
  assign ← ∅
  for each e ∈ G.F
     do f ← e
        repeat
              C ← ∅
              match ← MATCH-EQUATION(G, f, C, assign, vmap)
              if not match
                 then for each v ∈ C where v ∈ G.V
                          do let v′ be a vertex, such that v′ ∉ G.V
                             vmap[v] ← v′
                             G.V ← G.V ∪ {v′}
                      for each f ∈ C where f ∈ G.F
                          do let f′ be a vertex, such that f′ ∉ G.F
                             eqmap[f] ← f′
                             G.F ← G.F ∪ {f′}
                             for each v ∈ G.V where (f, v) ∈ G.E
                                 do G.E ← G.E ∪ {(f′, v), (f′, vmap[v])}
                      for each v ∈ C where v ∈ G.V
                          do assign[vmap[v]] ← eqmap[assign[v]]
                      f ← eqmap[f]
           until match
  return assign
```

Function Match-Equation returns TRUE. Consequently, we break out of the repeat-until loop and proceeds with the next equation.

Pendulum

```
MATCH-EQUATION(G, f, C, assign, vmap)
  C ← C ∪ {f}
  if there exits a v ∈ G.V such that (f, v) ∈ G.E
     and assign[v] = NIL and vmap[v] = NIL
    then assign[v] ← f
         return TRUE
    else for each v where (f, v) ∈ G.E and v ∉ C
           and vmap[v] = NIL
             do C ← C ∪ {v}
                if MATCH-EQUATION(G, assign[v], C, assign, vmap)
                  then assign[v] ← f
                       return TRUE
  return FALSE
```

Same pattern for the first four equations.



State after matching for f_1, f_2, f_3, f_4

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$$
$$eqmap = \{\}$$
$$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$$
$$C = \{f_4\}$$

Matched 4 equations (could also have matched lambda).

Note that only the last equation is colored because colors are cleared before matching.

Pendulum

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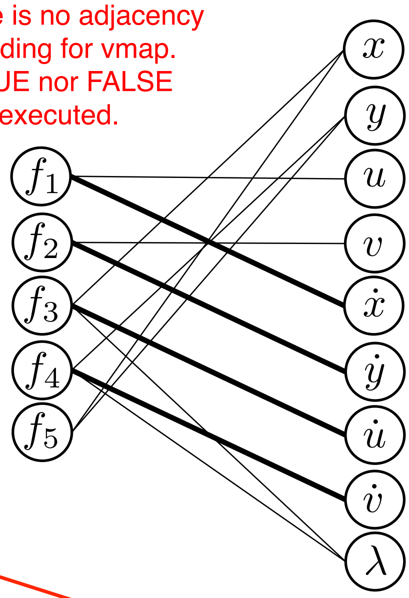
```
MATCH-EQUATION(G, f, C, assign, vmap)
  C ← C ∪ {f}
  if there exits a v ∈ G.V such that (f, v) ∈ G.E
     and assign[v] = NIL and vmap[v] = NIL
    then assign[v] ← f
         return TRUE
    else for each v where (f, v) ∈ G.E and v ∉ C
          and vmap[v] = NIL
           do C ← C ∪ {v}
              if MATCH-EQUATION(G, assign[v], C, assign, vmap)
                then assign[v] ← f
                     return TRUE
  return FALSE
```

For f_5, there is no adjacency variable holding for vmap. Neither TRUE nor FALSE branch are executed.



State after matching f_5

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$$
$$eqmap = \{\}$$
$$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$$
$$C = \{f_5\}$$

Algorithm Match-Equation returns FALSE and returns with f_5 colored.

Algorithm: Pantelides

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```
PANTELIDES(G, vmap, eqmap)
  assign ← ∅
  for each e ∈ G.F
    do f ← e
       repeat
             C ← ∅
             match ← MATCH-EQUATION(G, f, C, assign, vmap)
             if not match
               then for each v ∈ C where v ∈ G.V
                      do let v′ be a vertex, such that v′ ∉ G.V
                         vmap[v] ← v′
                         G.V ← G.V ∪ {v′}
                    for each f ∈ C where f ∈ G.F
                      do let f′ be a vertex, such that f′ ∉ G.F
                         eqmap[f] ← f′
                         G.F ← G.F ∪ {f′}
                         for each v ∈ G.V where (f, v) ∈ G.E
                           do G.E ← G.E ∪ {(f′, v), (f′, vmap[v])}
                    for each v ∈ C where v ∈ G.V
                      do assign[vmap[v]] ← eqmap[assign[v]]
                    f ← eqmap[f]
         until match
  return assign
```

We have match = FALSE

No colored variables.

But we have one colored equation.

No colored variables.

Pendulum

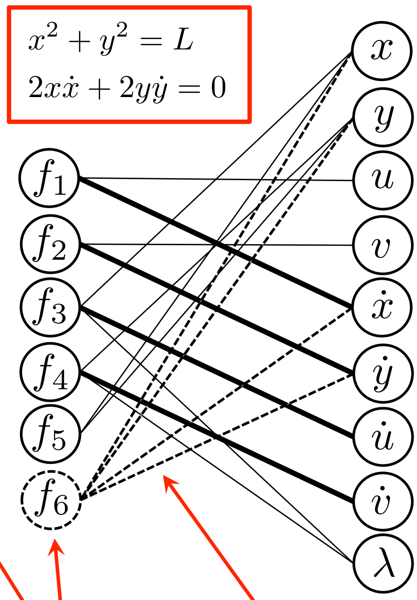
State after matching f_5

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$
$eqmap = \{\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_5\}$

for each $f \in C$ where $f \in G.F$
  do let f' be a vertex, such that $f' \notin G.F$
    $eqmap[f] \leftarrow f'$
    $G.F \leftarrow G.F \cup \{f'\}$
    for each $v \in G.V$ where $(f, v) \in G.E$
      do $G.E \leftarrow G.E \cup \{(f', v), (f', vmap[v])\}$

State after creating differentiated equation.

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$
$eqmap = \{f_5 \mapsto f_6\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_5\}$



Create a new equation node $\boldsymbol{f_6}$ by differentiating $\boldsymbol{f_5}$.

Create edges to variables and their derivatives.

Algorithm: Pantelides

PANTELIDES(G, $\underline{vmap}$, $\underline{eqmap}$)
$assign \leftarrow \emptyset$
for each $e \in G.F$
  do $f \leftarrow e$
    repeat
      $C \leftarrow \emptyset$
      $match \leftarrow$ MATCH-EQUATION($G, f, \underline{C}, \underline{assign}, vmap$)
      if not $match$
        then for each $v \in C$ where $v \in G.V$
          do let v' be a vertex, such that $v' \notin G.V$
            $vmap[v] \leftarrow v'$
            $G.V \leftarrow G.V \cup \{v'\}$
        for each $f \in C$ where $f \in G.F$
          do let f' be a vertex, such that $f' \notin G.F$
            $eqmap[f] \leftarrow f'$
            $G.F \leftarrow G.F \cup \{f'\}$
            for each $v \in G.V$ where $(f, v) \in G.E$
              do $G.E \leftarrow G.E \cup \{(f', v), (f', vmap[v])\}$
        for each $v \in C$ where $v \in G.V$
          do $assign[vmap[v]] \leftarrow eqmap[assign[v]]$
        $f \leftarrow eqmap[f]$
    until $match$
return $assign$

Repeat again (match was FALSE), but now with the differentiated equation $\boldsymbol{f_6}$.

$eqmap = \{f_5 \mapsto f_6\}$

Pendulum

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State before calling Match-Equation

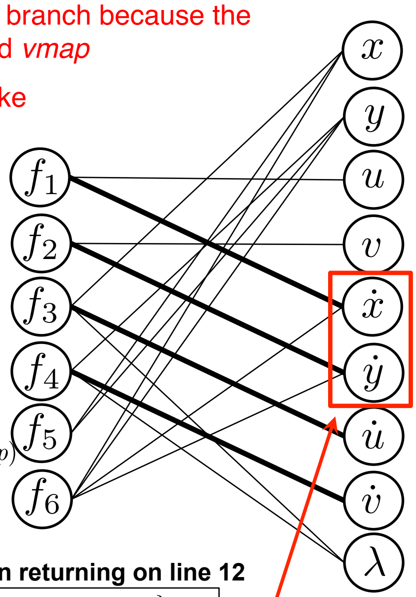
$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$
$eqmap = \{f_5 \mapsto f_6\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{\}$

Cannot take TRUE branch because the states of *assign* and *vmap*

Color and make recursive call.

```
MATCH-EQUATION(G, f, C, assign, vmap)
1  C ← C ∪ {f}
2  if there exits a v ∈ G.V such that (f, v) ∈ G.E
3      and assign[v] = NIL and vmap[v] = NIL
4     then assign[v] ← f
5          return TRUE
6     else  for each v where (f, v) ∈ G.E and v ∉ C
7           and vmap[v] = NIL
8              do C ← C ∪ {v}
9                 if MATCH-EQUATION(G, assign[v], C, assign, vmap)
10                  then assign[v] ← f
11                       return TRUE
12 return FALSE
```

Returns false.



Two variables are applicable for the FALSE branch.

Before first recursive call at line 9

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$
$eqmap = \{f_5 \mapsto f_6\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_6, \dot{x}\}$

After all recursion when returning on line 12

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$
$eqmap = \{f_5 \mapsto f_6\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$

Algorithm: Pantelides

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```
PANTELIDES(G, vmap, eqmap)
assign ← ∅
for each e ∈ G.F
   do f ← e
      repeat
           C ← ∅
           match ← MATCH-EQUATION(G, f, C, assign, vmap)
           if not match
             then for each v ∈ C where v ∈ G.V
                     do let v′ be a vertex, such that v′ ∉ G.V
                        vmap[v] ← v′
                        G.V ← G.V ∪ {v′}
                  for each f ∈ C where f ∈ G.F
                     do let f′ be a vertex, such that f′ ∉ G.F
                        eqmap[f] ← f′
                        G.F ← G.F ∪ {f′}
                        for each v ∈ G.V where (f, v) ∈ G.E
                           do G.E ← G.E ∪ {(f′, v), (f′, vmap[v])}
                  for each v ∈ C where v ∈ G.V
                     do assign[vmap[v]] ← eqmap[assign[v]]
                  f ← eqmap[f]
      until match
return assign
```

$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$

Differentiating equation two times

$$x^2 + y^2 = L$$
$$2x\dot{x} + 2y\dot{y} = 0$$
$$2x\ddot{x} + 2\dot{x}^2 + 2y\ddot{y} + 2\dot{y}^2 = 0$$

First step: create new differentiated variables

Pendulum

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State before creating new variables

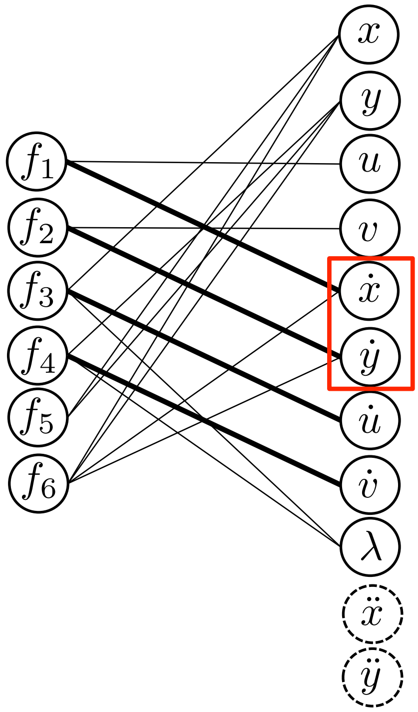
$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}\}$$
$$eqmap = \{f_5 \mapsto f_6\}$$
$$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$$
$$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$$

for each $v \in C$ where $v \in G.V$
  do let v' be a vertex, such that $v' \notin G.V$
    $vmap[v] \leftarrow v'$
    $G.V \leftarrow G.V \cup \{v'\}$

New variables and mapping

After adding new variables

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$$
$$eqmap = \{f_5 \mapsto f_6\}$$
$$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$$
$$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$$



Algorithm: Pantelides

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```
PANTELIDES(G, vmap, eqmap)
  assign ← ∅
  for each e ∈ G.F
      do f ← e
         repeat
               C ← ∅
               match ← MATCH-EQUATION(G, f, C, assign, vmap)
               if not match
                 then for each v ∈ C where v ∈ G.V
                          do let v' be a vertex, such that v' ∉ G.V
                             vmap[v] ← v'
                             G.V ← G.V ∪ {v'}
                      for each f ∈ C where f ∈ G.F
                          do let f' be a vertex, such that f' ∉ G.F
                             eqmap[f] ← f'
                             G.F ← G.F ∪ {f'}
                             for each v ∈ G.V where (f, v) ∈ G.E
                                 do G.E ← G.E ∪ {(f', v), (f', vmap[v])}
                      for each v ∈ C where v ∈ G.V
                          do assign[vmap[v]] ← eqmap[assign[v]]
                      f ← eqmap[f]
           until match
  return assign
```

Second step: create new differentiated equation nodes

Pendulum

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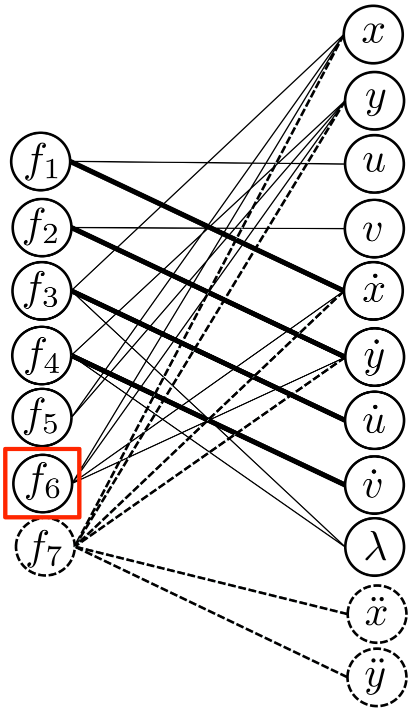
State before creating equation nodes

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$
$eqmap = \{f_5 \mapsto f_6\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$

for each $f \in C$ where $f \in G.F$
  do let f' be a vertex, such that $f' \notin G.F$
    $eqmap[f] \leftarrow f'$
    $G.F \leftarrow G.F \cup \{f'\}$
    for each $v \in G.V$ where $(f, v) \in G.E$
      do $G.E \leftarrow G.E \cup \{(f', v), (f', vmap[v])\}$

After adding equation f_6

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$
$eqmap = \{f_5 \mapsto f_6, f_6 \mapsto f_7\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$



Pendulum

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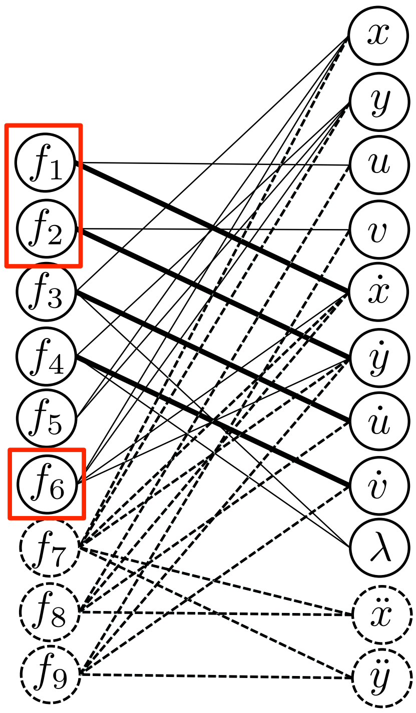
State before creating equation nodes

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$
$eqmap = \{f_5 \mapsto f_6\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$

for each $f \in C$ where $f \in G.F$
  do let f' be a vertex, such that $f' \notin G.F$
    $eqmap[f] \leftarrow f'$
    $G.F \leftarrow G.F \cup \{f'\}$
    for each $v \in G.V$ where $(f, v) \in G.E$
      do $G.E \leftarrow G.E \cup \{(f', v), (f', vmap[v])\}$

After adding all equations

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$
$eqmap = \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$
$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$



Algorithm: Pantelides

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```
PANTELIDES(G, vmap, eqmap)
assign ← ∅
for each e ∈ G.F
    do f ← e
       repeat
              C ← ∅
              match ← MATCH-EQUATION(G, f, C, assign, vmap)
              if not match
                then for each v ∈ C where v ∈ G.V
                          do let v′ be a vertex, such that v′ ∉ G.V
                             vmap[v] ← v′
                             G.V ← G.V ∪ {v′}
                     for each f ∈ C where f ∈ G.F
                          do let f′ be a vertex, such that f′ ∉ G.F
                             eqmap[f] ← f′
                             G.F ← G.F ∪ {f′}
                             for each v ∈ G.V where (f, v) ∈ G.E
                                 do G.E ← G.E ∪ {(f′, v), (f′, vmap[v])}
                     for each v ∈ C where v ∈ G.V
                          do assign[vmap[v]] ← eqmap[assign[v]]
                     f ← eqmap[f]
         until match
return assign
```

Third step: assign variables to equations for new variables.

Pendulum

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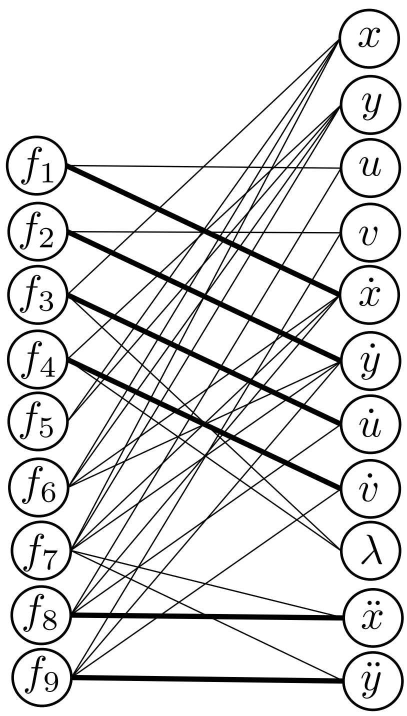
After adding all equations

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$$
$$eqmap = \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\}$$
$$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4\}$$
$$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$$

for each $v \in C$ where $v \in G.V$
  do $assign[vmap[v]] \leftarrow eqmap[assign[v]]$

After adding new assignments

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$$
$$eqmap = \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\}$$
$$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4, \ddot{x} \mapsto f_8, \ddot{y} \mapsto f_9\}$$
$$C = \{f_6, \dot{x}, f_1, \dot{y}, f_2\}$$



Algorithm: Pantelides

PANTELIDES(G, $\underline{vmap}$, $\underline{eqmap}$)

$assign \leftarrow \emptyset$
for each $e \in G.F$
 do $f \leftarrow e$
  repeat
   $C \leftarrow \emptyset$
   $match \leftarrow$ MATCH-EQUATION($G, f, \underline{C}, \underline{assign}, vmap$)
   if not $match$
    then for each $v \in C$ where $v \in G.V$
     do let v' be a vertex, such that $v' \notin G.V$
      $vmap[v] \leftarrow v'$
      $G.V \leftarrow G.V \cup \{v'\}$
    for each $f \in C$ where $f \in G.F$
     do let f' be a vertex, such that $f' \notin G.F$
      $eqmap[f] \leftarrow f'$
      $G.F \leftarrow G.F \cup \{f'\}$
      for each $v \in G.V$ where $(f, v) \in G.E$
       do $G.E \leftarrow G.E \cup \{(f', v), (f', vmap[v])\}$
    for each $v \in C$ where $v \in G.V$
     do $assign[vmap[v]] \leftarrow eqmap[assign[v]]$
    $f \leftarrow eqmap[f]$
  until $match$
return $assign$

Repeat again with second differentiated version of equation five.

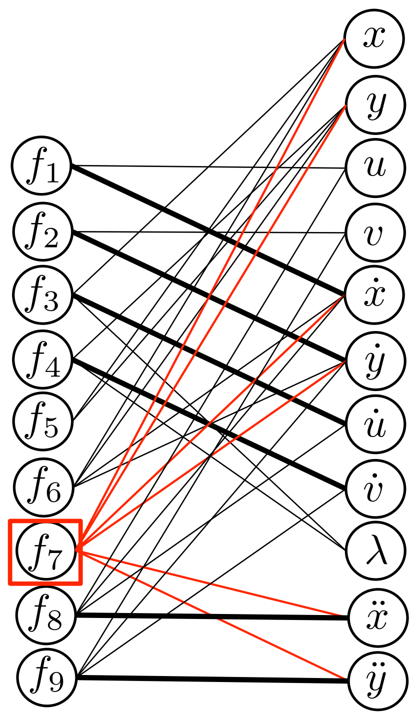
Pendulum

Before matching

$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$
$eqmap = \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\}$
$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4, \ddot{x} \mapsto f_8, \ddot{y} \mapsto f_9\}$

MATCH-EQUATION(G, f, $\underline{C}$, $\underline{assign}$, $vmap$)

$C \leftarrow C \cup \{f\}$
if there exits a $v \in G.V$ such that $(f, v) \in G.E$
 and $assign[v] =$ NIL and $vmap[v] =$ NIL
 then $assign[v] \leftarrow f$
   return TRUE
 else **for** each v where $(f, v) \in G.E$ and $v \notin C$
    and $vmap[v] =$ NIL
    do $C \leftarrow C \cup \{v\}$
     if MATCH-EQUATION(G, $assign[v]$, $\underline{C}$, $\underline{assign}$, $vmap$)
      then $assign[v] \leftarrow f$
        return TRUE
return FALSE



For clarity: view variables and edges where vmap[v] = NIL

Pendulum

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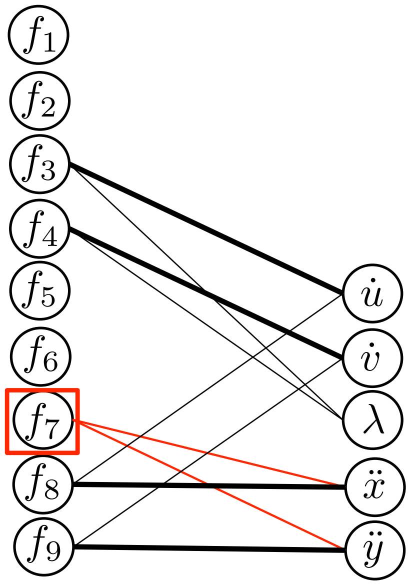
Before matching

$$
\begin{aligned}
vmap &= \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\} \\
eqmap &= \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\} \\
assign &= \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4, \\
&\quad\ \ddot{x} \mapsto f_8, \ddot{y} \mapsto f_9\}
\end{aligned}
$$

MATCH-EQUATION$(G, f, \underline{C}, \underline{assign}, vmap)$

$C \leftarrow C \cup \{f\}$
if there exits a $v \in G.V$ such that $(f, v) \in G.E$
  and $assign[v] =$ NIL and $vmap[v] =$ NIL
  then $assign[v] \leftarrow f$
    return TRUE
  else **for** each v where $(f, v) \in G.E$ and $v \notin C$
    and $vmap[v] =$ NIL
      do $C \leftarrow C \cup \{v\}$
        if MATCH-EQUATION$(G, assign[v], \underline{C}, \underline{assign}, vmap)$
          then $assign[v] \leftarrow f$
            return TRUE
return FALSE



Pendulum

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Before matching

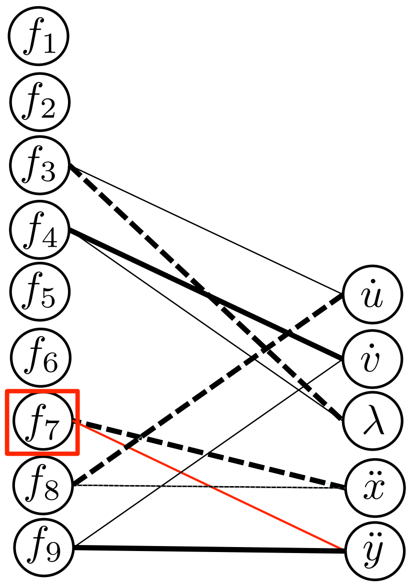
$$
\begin{aligned}
vmap &= \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\} \\
eqmap &= \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\} \\
assign &= \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_3, \dot{v} \mapsto f_4, \\
&\quad\ \ddot{x} \mapsto f_8, \ddot{y} \mapsto f_9\}
\end{aligned}
$$

Successful match!

MATCH-EQUATION$(G, f, \underline{C}, \underline{assign}, vmap)$

$C \leftarrow C \cup \{f\}$
if there exits a $v \in G.V$ such that $(f, v) \in G.E$
  and $assign[v] =$ NIL and $vmap[v] =$ NIL
  then $assign[v] \leftarrow f$
    return TRUE
  else **for** each v where $(f, v) \in G.E$ and $v \notin C$
    and $vmap[v] =$ NIL
      do $C \leftarrow C \cup \{v\}$
        if MATCH-EQUATION$(G, assign[v], \underline{C}, \underline{assign}, vmap)$
          then $assign[v] \leftarrow f$
            return TRUE
return FALSE

$$
\begin{aligned}
vmap &= \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\} \\
eqmap &= \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\} \\
assign &= \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_8, \dot{v} \mapsto f_4, \\
&\quad\ \lambda \mapsto f_3, \ddot{x} \mapsto f_7, \ddot{y} \mapsto f_9\} \\
C &= \{f_7, \ddot{x}, f_8, \dot{u}, f_3\}
\end{aligned}
$$



Algorithm: Pantelides

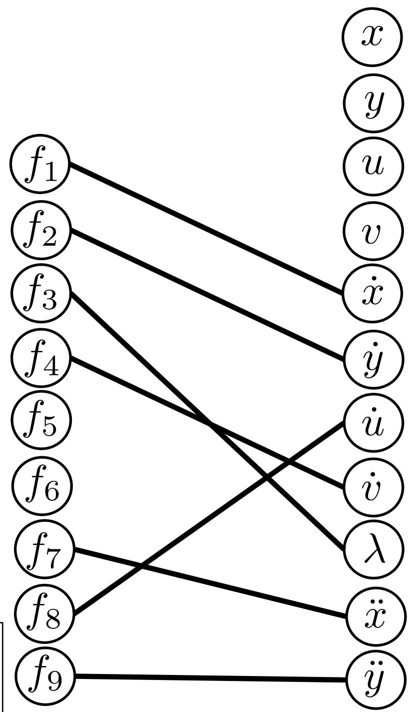
```
PANTELIDES(G, vmap, eqmap)
assign ← ∅
for each e ∈ G.F
    do f ← e
        repeat
            C ← ∅
            match ← MATCH-EQUATION(G, f, C, assign, vmap)
            if not match
              then for each v ∈ C where v ∈ G.V
                       do let v′ be a vertex, such that v′ ∉ G.V
                          vmap[v] ← v′
                          G.V ← G.V ∪ {v′}
                   for each f ∈ C where f ∈ G.F
                       do let f′ be a vertex, such that f′ ∉ G.F
                          eqmap[f] ← f′
                          G.F ← G.F ∪ {f′}
                          for each v ∈ G.V where (f, v) ∈ G.E
                              do G.E ← G.E ∪ {(f′, v), (f′, vmap[v])}
                   for each v ∈ C where v ∈ G.V
                       do assign[vmap[v]] ← eqmap[assign[v]]
                   f ← eqmap[f]
        until match
return assign
```

Last equation and successful match.
Algorithm terminates.

Result of Pantelides on Pendulum

$$vmap = \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\}$$
$$eqmap = \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\}$$
$$assign = \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_8, \dot{v} \mapsto f_4, \lambda \mapsto f_3, \ddot{x} \mapsto f_7, \ddot{y} \mapsto f_9\}$$

(1)	$\dot{x} = u$	$f_1(\dot{x}, u) = 0$
(2)	$\dot{y} = v$	$f_2(\dot{y}, v) = 0$
(3)	$\dot{u} = \lambda \cdot x$	$f_3(\dot{u}, \lambda, x) = 0$
(4)	$\dot{v} = \lambda \cdot y - g$	$f_4(\dot{v}, \lambda, y) = 0$
(5)	$x^2 + y^2 = L$	$f_5(x, y) = 0$
(6)	$2x\dot{x} + 2y\dot{y} = 0$	$f_6(x, \dot{x}, y, \dot{y}) = 0$
(7)	$2x\ddot{x} + 2\dot{x}^2 + 2y\ddot{y} + 2\dot{y}^2 = 0$	$f_7(x, \dot{x}, \ddot{x}, y, \dot{y}, \ddot{y}) = 0$
(8)	$\ddot{x} = \dot{u}$	$f_8(\ddot{x}, \dot{u}) = 0$
(9)	$\ddot{y} = \dot{v}$	$f_9(\ddot{y}, \dot{v}) = 0$



$|G.F| = 9 \qquad |G.V| = 11$

Two variables out of the set G.V can be given arbitrary initialization values, as long as all constraints above are satisfied.

Result of Pantelides on Pendulum

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$$
\begin{aligned}
vmap &= \{x \mapsto \dot{x}, y \mapsto \dot{y}, u \mapsto \dot{u}, v \mapsto \dot{v}, \dot{x} \mapsto \ddot{x}, \dot{y} \mapsto \ddot{y}\} \\
eqmap &= \{f_5 \mapsto f_6, f_6 \mapsto f_7, f_1 \mapsto f_8, f_2 \mapsto f_9\} \\
assign &= \{\dot{x} \mapsto f_1, \dot{y} \mapsto f_2, \dot{u} \mapsto f_8, \dot{v} \mapsto f_4, \\
&\qquad \lambda \mapsto f_3, \ddot{x} \mapsto f_7, \ddot{y} \mapsto f_9\}
\end{aligned}
$$

Is the system of equations solvable if we replace the old equations with their differentiated version?

$$
\begin{aligned}
&(1)\quad \dot{x} = u \\
&(2)\quad \dot{y} = v \\
&(3)\quad \dot{u} = \lambda \cdot x \\
&(4)\quad \dot{v} = \lambda \cdot y - g \\
&(5)\quad x^2 + y^2 = L \\
&(6)\quad 2x\dot{x} + 2y\dot{y} = 0 \\
&(7)\quad 2x\ddot{x} + 2\dot{x}^2 + 2y\ddot{y} + 2\dot{y}^2 = 0 \\
&(8)\quad \ddot{x} = \dot{u} \\
&(9)\quad \ddot{y} = \dot{v}
\end{aligned}
$$

By substituting (8) and (9) we have

$$
\begin{aligned}
\ddot{x} &= \lambda \cdot x \\
\ddot{y} &= \lambda \cdot y - g \\
2x\ddot{x} + 2\dot{x}^2 &+ 2y\ddot{y} + 2\dot{y}^2 = 0
\end{aligned}
$$

Which is solvable for highest derivative

	λ	$\ddot{y}$	$\ddot{x}$
f_1	**1**	0	**1**
f_2	**1**	**1**	0
f_3	0	**1**	**1**

Same result if converted into order one equation

	$\dot{y}$	$\dot{x}$	$\dot{u}$	λ	$\dot{v}$
f_2	**1**	0	0	0	0
f_1	0	**1**	0	0	0
f_5	**1**	**1**	**1**	0	**1**
f_3	0	0	**1**	**1**	0
f_4	0	0	0	**1**	**1**

Termination of Pantelides

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Does Pantelides algorithm terminate?

Depends on the input graph.

Before Pantelides, check that matching can be found on a matrix where variables do not distinguish if they appear differentiated or not.

$$
\begin{aligned}
\dot{x} &= u \\
\dot{y} &= v \\
\dot{u} &= \lambda \cdot x \\
\dot{v} &= \lambda \cdot y - g \\
x^2 + y^2 &= L
\end{aligned}
$$

	x	y	u	v	λ
f_1	**1**	0	**1**	0	0
f_2	0	**1**	0	**1**	0
f_3	**1**	0	**1**	0	**1**
f_4	0	**1**	0	**1**	**1**
f_5	**1**	**1**	0	0	0

Matrix to check

	u	λ	v	y	x
f_1	**1**	0	0	0	**1**
f_3	**1**	**1**	0	0	**1**
f_4	0	**1**	**1**	**1**	0
f_2	0	0	**1**	**1**	0
f_5	0	0	0	**1**	**1**

Yes, match was found. Hence the problem is not structurally singular.

Part IV

Dummy Derivatives

Index Reduction

Should differentiated equations from Pantelides be used for index reduction?

$$\ddot{x} = \lambda \cdot x$$
$$\ddot{y} = \lambda \cdot y - g$$
$$2x\ddot{x} + 2\dot{x}^2 + 2y\ddot{y} + 2\dot{y}^2 = 0$$

The reduced problem (index-1) is mathematically correct, but since equation

$$x^2 + y^2 = L$$

is not present, numerical approximation gives a "drifting problem". In our example, the pendulum's length will grow...

Dummy Derivative

Basic Idea:

- Include all differentiated equations
- For each equation, introduce a "dummy derivative" variable.

$$\ddot{x} = \lambda \cdot x$$
$$y'' = \lambda \cdot y - g$$
$$x^2 + y^2 = L$$
$$2x\dot{x} + 2yy' = 0$$
$$2x\ddot{x} + 2\dot{x}^2 + 2yy'' + 2y'^2 = 0$$

All constraints are present and the number of equations and unknowns match.

The actual algorithm is presented by Mattson and Söderlind (1993)

References and Further Reading

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- C. C. Pantelides. The Consistent Initialization of Differential-Algebraic Sys- tems. SIAM Journal on Scientific and Statistical Computing, 9(2):213–231, 1988.
- R. Tarjan. Depth-first search and linear graph algorithms. SIAM Journal on Computing, 1(2):146–160, 1972.

Conclusions

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Matching
Finds a mapping between variables and equations. Used both in BLT sorting and Pantelides algorithm

BLT Sorting
Sort blocks of equation, where each block represents an algebraic loop. Uses matching and Tarjan's algorithm

Pantelides
Determine the subset of equations that needs to be differentiated.

Dummy Derivative
Method that uses Pantelides to perform correct index reduction.

Thank you for listening!